

Topic: Logic programming

Lectures: Monday August 11, Tuesday August 12

Reading: Abelson & Sussman, Section 4.4.1–3

This assignment gives you practice writing logic programs. It's very much "above the line" since we don't expect you to know how the query system works. This homework is due at **midnight on Wednesday, August 13**. Please put your solutions into a file `hw8-1.scm` and submit electronically. Make sure to include your test cases, too.

To add an assertion: `(assert! <conclusion>)`

To add a rule: `(assert! (rule <conclusion> <body (optional)>))`

Anything else is a query.

The query interpreter is in the file `~cs61a/lib/query.scm`. To initialize the interpreter type `(query)`; to re-enter the main loop without reinitializing, type `(query-driver-loop)`. Nothing—not even the rules for the `same` and `append-to-form` relations—is there when the interpreter is initialized.

Question 1. Do Exercise 4.56 in SICP. To load the database, type the following after loading `query.scm`:

```
STk> (initialize-data-base microshaft-data-base)
STk> (query-driver-loop)
```

The pattern `(?a . ?b)` matches any pair, so you can use it to print everything that is in the database.

Question 2. This question explores the unary arithmetic system described in lecture where numbers are represented as lists.

- A.** Note that summing two of these unary numbers merely involves joining the lists that represent them. We can define a rule for adding query numbers using `append-to-form` (Page 451):

```
;;; Query input:
(assert! (rule (?a + ?b = ?c) (append-to-form ?a ?b ?c)))
Assertion added to data base.
;;; Query input:
((a a a a) + (a a a) = ?what)
;;; Query results:
((a a a a) + (a a a) = (a a a a a a a)) ;; 4 + 3 = 7
```

Devise rules to allow multiplication of query numbers:

```
;;; Query input:
((a a a a) * (a a a) = ?what)
;;; Query results:
((a a a a) * (a a a) = (a a a a a a a a a a a)) ;; 4 * 3 = 12
```

The question continues on the next page.

B. Using your multiplication rule from above, implement a factorial relation for query numbers:

```
;;; Query input:
((a a a a) != ?what)
;;; Query output:
((a a a a) != (a a a a a a a a a a a a a a a a a a a a a a)) ;; 4! = 24
```

Question 3. We can interpret the query interpreter's failure to return any results as saying, "Your query was not consistent with any assertions I know or any rules I can apply on the basis of those assertions." This can be used to implement true/false queries where the interpreter echoes the query if it is true and displays no results otherwise. For example:

```
;;; Query input:
(deep-member a ((a) b))
;;; Query output:
(deep-member a ((a) b)) ;; true
;;; Query input:
(deep-member c ((b ((a))))))
;;; Query output:
;; false
;;; Query input:
(deep-member c ((b ((a c))))))
;;; Query output:
(deep-member c ((b ((a c))))))
;;; Query input:
(deep-member ?anything ())
;;; Query output:
;; false
```

Write rules for the `deep-member` relation that behaves as above.

Question 4. Write query rules for the `assoc` relation. It should work like this:

```
;;; Query input:
(assoc carolen ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) ?what)
;;; Query results:
(assoc carolen ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) (carolen 10))
;;; Query input:
(assoc todd ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) ?what)
;;; Query results:
;; no results!
```

Notice that the *first* sublist beginning with `carolen` is brought forth. The query should run backward, too:

```
;;; Query input:
(assoc ?who ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) (?who 10))
;;; Query results:
(assoc greg ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) (greg 10))
(assoc carolen ((greg 10) (kurt 12) (carolen 10) (alex 13) (carolen 15)) (carolen 10))
```